# Overview

## Focus on regression problem

Which variables ought to be used in a regression, particularly when the number of potential predictors p is large (data mining).

## **Reproduce criteria**

Model selection via *AIC*, *BIC*, *RIC*, *eBIC* are equivalent to choosing the model which offers the greatest compression according to a specific two-part code. Selection criteria differ in how they encode the parameters.

# Composition of prefix

Prefix must indicate two features:

- 1. Coefficient estimates
- 2. Which variables are being used

### Similar to location problem

Again will encode parameters as integer z-scores, adding enough information to associate estimates with variables.

# **Regression Model**

### True Model

Rather than assume  $EY = X\beta$ , mean is unspecified:

$$Y = \eta + \epsilon, \quad E \epsilon = 0, \quad \operatorname{Var} \epsilon = \sigma^2 I_n,$$

### Working Model

View  $X\beta$  as the projection of  $\eta$  into the space defined by the available set of predictors,

 $Y = X\beta + \epsilon$  where  $X\beta = (X(X'X)^{-1}X)\eta$ ,

and treat  $\epsilon \sim N(0, \sigma^2)$ .

#### Subset/selection coefficients

Let  $\gamma = (\gamma_1, \dots, \gamma_p)$  denote a sequence of p 0's and 1s. Denote a subset of  $\beta$  by (miss APL notation!)

$$\beta_{\gamma}$$
 defined by  $\beta_j \neq 0 \iff \gamma_j = 1$ 

#### Simplifying assumptions

- $p \leq n$  possible orthogonal regressors  $X_j$ , with  $X'_j X_j = n$ .
- $\sigma^2$  is known.
- Receiver knows n and X, so needn't send either.

# **Coding Regression**

### Model prefix

Prefix encodes  $\gamma$  and associated estimates:

- 1. Code for  $\gamma = (1, 0, 0, \dots, \gamma_j, \dots, 1)$ , the selection indicator
- 2. Code for fitted  $\hat{\beta}_{\gamma}$  estimates
- 3. Compressed data

### Goal and protection

The goal remains to construct the shortest message. Note the automatic penalty for over-fitting: the more variables used, the longer the prefix since more estimates must be added.

### Estimates

$$b_j = \beta_j + \frac{X'_j \epsilon}{X'_j X_j} = \beta_j + \frac{\sigma}{\sqrt{n}} Z, \quad Z \sim N(0, 1)$$

so that  $SE(b_j) = \sigma/\sqrt{n}$ 

#### **Rounding coefficients**

Round coefficient estimates  $b_j$  to a standard error scale, as in the location problem,

$$\hat{\beta}_j = \frac{\sigma \langle z_j \rangle}{\sqrt{n}} , \qquad z_j = \frac{\sqrt{n} \, b_j}{\sigma}$$

# Variable Selection via Coding

### Trade-off

Add additional variable  $X_j$  if

(Gain in data compression) > (Increase in prefix length).

#### Gain in data compression

Log-likelihood based on q predictors is (ignoring constants)

$$\log \frac{1}{P(Y|b_{\gamma_q})} = \frac{\sum_{i=1}^n (Y_i - \hat{Y}_i(q))^2}{\sigma^2 2 \ln 2} = \frac{RSS(q)}{\sigma^2 2 \ln 2}$$

If add another predictor, say  $X_j$ , then

$$\triangle RSS = RSS(q) - RSS(q+1) = n \, b_j^2 \; ,$$

so the gain in data compression is

$$\frac{\triangle RSS}{\sigma^2 2 \ln 2} = \frac{nb_j^2}{\sigma^2 2 \ln 2} = \frac{z_j^2}{2 \ln 2}$$
 fewer bits.

#### Parameter cost differs

Least squares, AIC, BIC, RIC, and eBIC code the parameters

 $\gamma, \hat{eta}$ 

differently, and so reach different compromises of data compression and model complexity.

# Least Squares Coding

### Fixed format code

Fixed format with reserved, fixed-length fields for each predictor:

- 1. p bits for the indicator  $\gamma$ ,
- 2.  $\frac{k}{2}\log n$  for each parameter  $(2|\beta_j| < k\sigma/\sqrt{n})$ , and
- 3. tag on the fully compressed data.

## Parameters of code

$$\underbrace{\frac{\gamma_1, \dots, \gamma_p}{p \text{ bits}}}_{\text{ all } p} \| \underbrace{\frac{\langle z_1 \rangle}{(k/2) \log n \text{ bits}}}_{(k/2) \log n \text{ bits}} \underbrace{\frac{\langle z_p \rangle}{(k/2) \log n \text{ bits}}}_{\text{ all } p}$$

#### Data

Encode data using the associated rounded parameter estimates. This requires about

$$\log \frac{1}{P(Y_1, \dots, Y_n | b_1, \dots, b_p)} + \underbrace{nQ}_{\text{quantized}}$$
 bits

### **Resulting selection**

With fixed-length fields, regardless of selected variables, one obtains the shortest message by encoding *all* of the parameters.

# BIC/SIC Code

## Partly-fixed format code

Simple modification of the OLS code, with a fixed format for each *chosen* predictor rather than all predictors:

- 1. p bits for the indicator  $\gamma$ ,
- 2.  $\frac{k}{2} \log n$  for each chosen parameter, and
- 3. tag on the fully compressed data.

**Parameters of code** Assuming  $|\gamma| = \sum_j \gamma_j = q$ ,



## **Resulting selection**

Add  $X_j$  if



Increased compression

Increased parm bits

implying that one selects  $X_j$  if (with k = 1)

$$|z_j| > \sqrt{\log n} \; ,$$

as when using *BIC*.

# Interpreting the BIC/SIC Code

#### Code

Assuming  $|\gamma| = \sum_j \gamma_j = q$ ,



#### Spike and slab for each slope

When  $X_j$  is

**Excluded:** 1 bit to denote zero (in the code for  $\gamma_j$ ).

**Included:**  $1 + \frac{1}{2} \log n$  bits for  $\gamma_i$  and  $z_j$ .

#### Prior on "complexity"

Since  $\gamma$  is always coded in p bits, as though *iid* coin tosses, this code assigns equal probability to all  $2^p$  possible models.

Prior 
$$\operatorname{prob}(q=0) = \frac{1}{2^p}$$
, Prior  $\operatorname{prob}(q=1) = \frac{p}{2^p}$ 

In general

Prior 
$$\operatorname{prob}(q) = \frac{\binom{p}{q}}{2^p}$$

so that the most favored model (highest prior) is q = p/2.  $\Rightarrow$  We expect *half* of the variables to enter the model, albeit with a high threshold,  $|z_j| > \sqrt{\log n}$ .

# **AIC Regression Coding**

## Variable length code

Fixed *p*-bit prefix for  $\gamma$  with varying fields for each predictor:

- Prefix  $\gamma$  embedded in parameter codes,
- Concatenate universal codes  $U_s(z_j), j = 1, \ldots, p$  for  $z_j$ .

### Parameters of code

$$\underbrace{U_s(z_1), \ U_s(z_2), \ \ldots, U_s(z_p)}_{\cdot \cdot \cdot \cdot \cdot}$$

p univ codes

EG: p = 6 and simpler to read Cauchy codes,

**0 1** 0 + **0 1** 1 1 0 - **0**  $\Rightarrow$  0 1 0 - 3 0 0

Leading bits of the universal codes are indicators  $\gamma_j$ .

#### **Resulting selection**

Add  $X_j$  if improved goodness of fit compensates for adding  $U_s(z_j)$  bits for the additional parameter,



which implies that one codes once  $|z_j| > 2.4$  (approximately).

#### **Resembles AIC**

Threshold fixed on z scale, as with AIC or  $C_p$ .

# Interpreting the AIC Code

## Associated prior on $\beta_j$

The associated prior on coordinates is a "rectangular"

log-Cauchy distribution, and is *not* spherically symmetric.

Coefficients are not artificially constrained to some interval.

### Natural prior?

Suggests many small coefficients, regardless of the sample size.

## Prior on complexity

Though embedded into universal codes,  $\gamma$  still uses p bits for all models, as in the *BIC* code; expect half of the variables to enter.

### Local asymptotics

Motivated by local asymptotics in which one fixes z as  $n \to \infty$ rather than letting the z score grow to infinity.

# **Robustness of Universal Codes**

### What about the message lengths?

If in fact the z scores are large, won't the AIC model codes be a lot longer than the BIC model codes?

### Oracle

Suppose that it is known that  $-M\sigma \leq \beta_j \leq M\sigma$ .

### Data

Suppose further that the fitted coefficients for q variables attain this upper limit,  $b_1 = \cdots = b_q = M\sigma$  so that  $z_j = \sqrt{n}M$ .

Large value exceeds AIC and BIC thresholds.

**Prefix length for BIC code** Since the grid of z-scores has

$$\frac{2M\sigma}{\sigma/\sqrt{n}} = 2\sqrt{n}\,M \quad \text{positions},$$

uniform coding requires

$$q \log 2\sqrt{n}M = q(\log \sqrt{n}M + 1)$$
 bits.

Prefix length for AIC code

$$q \log^* \sqrt{n}M \approx q \left(\log \sqrt{n}M + 2\log\log \sqrt{n}M\right)$$

which shares the dominant term with the BIC code.

# MDL in Regression

**Definition of minimum description length**Rissanen (1983)In its original asymptotic form, the MDL for a model with qorthogonal predictors is

$$MDL(q) = \log^*(V(k) \|\hat{\theta}\|^q) + \log \frac{1}{P(Y|\hat{\theta})}$$
  

$$\approx \frac{q}{2} \log n + \frac{q}{2} \log \sum \hat{\theta}_i^2 + \log V(q) + \log \frac{1}{P(Y|\hat{\theta})}$$
  

$$\approx \frac{q}{2} \log n + \log \frac{1}{P(Y|\hat{\theta})}$$
  

$$= BIC(q)$$

where V(k) = Vol(k-dim ball, radius one) and  $\|\theta\|^2 = n \sum \theta_j^2$ .

### Implicit selection indicator

Since  $\gamma$  does not appear in this definition, its as though it is coded with a fixed number of bits for all models.

### Local coding interpretation

The prefix encodes an index for vector  $\mathbf{z} = (z_1, \ldots, z_q)$  using a spiraling code:

$$MDL(q) \approx \log^*(V(q) \|\hat{\theta}\|^q) + \log 1/P(Y|\hat{\theta})$$
  
=  $U(i(\mathbf{z})) + \log 1/P(Y|\hat{\theta})$ 

where  $i(\mathbf{z})$  denotes index identifying the vector z score.

# Spiral MDL Index Path

- "Spiral" indexing using universal code on SE scale.
- Plots of index (below), bits for index (next).



# Code Lengths for Spiral MDL Index



# Implications of Spherical Prior/Code

### Projection to subspace

Following plot: code lengths with subspaces show that threshold for adding another variable *increases* with  $\|\mathbf{z}\|$ .

### Example: Use one or two variables?

q = 2 gives shorter code than q = 1 when

$$\overline{2} > \underbrace{U(i(z_1, z_2)) - U(i(z_1))}_{z_2}$$

Gain in compression

Increase parm bits

For large  $z_1 >> z_2 > 0$ ,

$$U(i(z_1, z_2)) - U(i(z_1)) \approx \log^* ||z_1, z_2||^2 - \log^* ||z_1||$$
  
$$\approx 2 \log z_1 - \log z_1$$
  
$$= \log z_1$$

Add  $X_2$  to model with just  $X_1$  when

$$\frac{z_2^2}{2\ln 2} > \log z_1$$

### Maximum coefficient determines threshold

q large coefficients,  $\mathbf{z}=(z,z,\ldots,z),$  and one smaller coefficient,  $z>>\tilde{z}>0$  . Add  $\tilde{z}$  if

$$\frac{\tilde{z}^2}{2\ln 2} > \log z + \frac{1}{2}\log q$$

# **Code Lengths with Projections**



# Examples

Add  $\tilde{z}$ ?

Add  $\tilde{z}$  to model with q coefficients when

$$\frac{\tilde{z}^2}{2\ln 2} > \underbrace{\log z + \frac{1}{2}\log q}_{\text{penalty}}$$

Local coding = BIC in special case

If max  $\mathbf{z} = \sqrt{n}$ , then penalty  $= \frac{1}{2}(\log n + \log q)$  and threshold is about  $\sqrt{\log n}$ .

Explicit examples "Elephants and mice"

Model A Two small coefficients

$$\mathbf{z} = (3,4) \Rightarrow \text{pick} \Rightarrow (3,4)$$

Model B Two small, plus one large

$$\mathbf{z} = (3, 4, 10) \Rightarrow \text{pick} \Rightarrow (3, 4, 10)$$

 $\mathbf{z} = (3, 4, 100) \quad \Rightarrow \text{ pick } \Rightarrow \quad (4, 100)$ 

 $\mathbf{z} = (3, 4, 1000) \implies \text{pick} \implies (1000)$ 

Model C Two small, plus many large

 $\mathbf{z} = (3, 4, 100, \dots, 100) \implies \text{pick} \implies (4, 100, \dots, 100)$ 

# Indexed Parameter Coding: RIC

# Complexity prior

Suppose expect very few coefficients to enter model,  $|\gamma| \approx 1$ .

## Bernoulli compression

Compress  $Y_1, \ldots, Y_n \sim B(1/n)$  by giving indices for *i* s.t.  $Y_i = 1$ ,

$$n H(Y_1) \approx \log n$$

Encode  $\gamma$  as a sequence of indices rather than 0/1 indicators.

Parameters of code

$$q \mid \underbrace{(j_1, U_s(z_{j_1}))}_{\log p + \ell(U_s(z_{j_1}))} \mid \cdots \mid (j_q, U_s(z_{j_q}))$$

### **Resulting selection**

Add  $X_j$  if (approximately)

$$\underbrace{z_j^2/(2\ln 2)}_{\text{Increased compression}} > \underbrace{\log p + 2\log \langle z_j \rangle}_{\text{Increased parm bits}}$$

or roughly once  $z_j$  exceeds the *Bonferroni* bound,

$$|z_j| > \sqrt{2 \log p} \approx \Phi^{-1}(1 - 1/p)$$

# Adaptive Coding: eBIC

# Coding methods

Method	Code for $\gamma$	Expected Predictors
BIC	p bit prefix	p/2
AIC	Embedded $p$ bit prefix	p/2
RIC	Indexing	1

Why a priori assume the complexity — code adaptively.

### Compress $\gamma$

Compress  $\gamma_1, \ldots, \gamma_p$ , treating as a Bernoulli sequence. In effect, modify the *AIC* code by compressing the leading bits of the sequence of universal codes.

Such a code will produce slightly longer messages than

- RIC code if indeed q = 1 is best
- AIC code if q = p/2 is best

but the added length in these cases is very small.

Two-part code

$$\underbrace{\gamma_1, \dots, \gamma_p}_{p \ H(q/p)} \mid \underbrace{U_s(z_{j_1}), \dots, U_s(z_{j_q})}_{\sum_{k=1}^q \ell(U(z_{j_k}-q))}$$

# Adaptive Coding: eBIC

# Adaptive selection criterion

Add  $X_j$  if (approximately)

$$\frac{z_j^2}{2\ln 2} > p(H(\frac{q+1}{p}) - H(\frac{q}{p})) + 2\log z_j$$

or once  $|z_j|$  exceeds the adaptive thresholding bound,

$$\frac{z_j^2}{2\ln 2} > \log \frac{p-q}{q+1} + 2\log z_j \quad \Rightarrow \quad |z_j| > \sqrt{2\log p/q}$$

### Comparable to

- Simes, Step-up/Step-down Testing: Compare  $\max z_j$  to Bonferroni, second largest to next normal order stat, etc.
- Empirical Bayes prior for the number of parameters.

## Further variations

Can compresses other bits in the z scores as well.

### **Big** question

Is coding useful when used in this more extensive manner?

- Fine for offering another way to think about model selection criteria.
- But, is it appropriate to use bit lengths to judge which is a better selection criterion?

# Discussion

## Model selection message formats

Criterion	Threshold	Parameter	Complexity
BIC, SIC	$\sqrt{\log n}$	Spike-slab	half
$AIC, C_p$	$\sqrt{2}$	Log-Cauchy	half
<i>RIC</i> , hard	$\sqrt{2\log p}$	Log-Cauchy	1
eBIC	$\sqrt{2\log p/q}$	Log-Cauchy	adaptive

## Discussion

- Interplay of information theory and Bayesian ideas Another way to think about priors, particularly in harder problems (priors for continuous functions).
- Robustness of the priors Universal priors are uncommon, but quite powerful.
- Spherical priors

Common, but appropriate in variable selection?

• Is coding a realistic criterion?

Seems fine as a way to characterize methods, but can it suggest which are really better?

# **Research Continues**

# Other types of models

Application to smoothing splines, piecewise models, or regression trees?

Varying parameter effects on compression, and how to code?

### **Dependent** processes

Context tree models are capable of capturing dependence in a nonparametric way and have been used to develop, for example, model-free bootstrap resampling methods.

# Reducing assumptions: Collinearity, variance, normality Collinearity Drop the assumption of orthogonal parameters.

Variance Drop the assumption of known variance.

**Robustness** CLT for coefficients, but what about compression of the data via likelihood? Using wrong distribution makes a big difference in *data compression*, implying  $z^2/(2 \ln 2)$  may not be the right trade-off.

### File compression vs model selection

Coding seeks big gains in compression, on the order of 10%. Testing/parameter selection deals with several bits.

Can a method used to obtain 10% gains be applied to evaluate changes of several bits?