Stochastic Combinatorial Optimization from TSPs to MSTs via Caterpillars and Dogerpillars

J. Michael Steele University of Pennsylvania, Wharton School

J. Michael Steele University of Pennsylvania, Wharton School Stochastic Combinatorial Optimization from TSPs to MSTs via

▶ Part 1: Getting Started with the MST and TSP

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- ▶ Part III: Interpolation The Real Theme
  - (Wherein the famous Dogerpillars are introduced and explored.)

#### Consider a set of points ...





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On the Other Hand:

There is a O(n) time  $\epsilon$  algorithm (Karp-Steele (1985)) if you assume a probability model for the points.

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The MST and TSP may LOOK like similar problems ....

► BUT:

Their computational theory tells us that they are wildly different

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For the TSP this is the famous Beardwood-Halton-Hammersly theorem of 1959. For the MST the result is from Steele (1988). The constants C<sub>TSP</sub> and C<sub>MST</sub> are not known exactly. The natural analogs hold in d ≥ 2.

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- The modern approach to the variance is radically different from that used by BHH
- The modern package is much more robust to changes in "problem" and "model".

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Even now this may seem surprising. Here, and in many other cases, it gives an very pleasing path to the desired strong laws.

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We do not know if this is the "truth" when d > 2. Lower bounds on variance are hard to come by. We have similar open issues with respect to sharp concentration in d > 2.

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- ► Three Hints Do It: Cyclic processes, O(n<sup>-(1+ϵ)/2</sup>) shifts, and subsequent scales.

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- The minimal spanning caterpillar is well defined and we naturally have

$$L_n^{\mathsf{MST}} \leq L_n^{\mathsf{CAT}} \leq L_n^{\mathsf{TSP}}$$

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Remove a path and have only stars ...



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# Strong Law for Caterpillars

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- One suspects this can be improved to universal boundedness as for the TSP and MST
- Still, this is good enough. One gets the strong law for minimal spanning caterpillars.

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- Ironically, if you Google "dogerpillar" you still find much that irrelevant. It's hard to come up with a neologism these days.

Introducing Dogerpillars

## Introducing Dogerpillars

Definition: A graph G is a dogerpillar (more precisely a k = k(n)-dogerpillar) if

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#### Specializations:

- Taking k = n, we see that the Dogerpillar is a tree
- Taking k = 0, we see that the Dogerpillar is a path
- The most interesting cases are k = o(n), especially k = O(√n).

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- Let's look at the progress to date...

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$$\lim_{n \to \infty} \frac{L_n^{\text{DOG}}}{\sqrt{n}} = C \int_{R^2} \sqrt{f(x)} dx$$

• For all  $k_n \sim n^{\alpha}$  with  $0 < \alpha < 1$  one has the strong law:

1

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  - CLT for at least some interesting ranges of  $k_n$ , at least  $k_n \sim n^{1/2+\epsilon}$

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We've reviewed the probability theory of the TSP and MST as it has evolved over the last 25 years.

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- Thank You for Your Attention!